

SUSE® Storage

Sizing and Performance for Ceph

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What is Ceph?

From 10,000 Meters

- Open Source Storage Distributed solution
- Most popular choice of distributed storage for openStack^[1]
- Lots of goodies
 - Distributed Object Storage
 - Redundancy
 - Efficient Scale-Out
 - Can be built on commodity hardware
 - Lower operational cost

[1] <http://www.openstack.org/blog/2013/11/openstack-user-survey-statistics-november-2013/>

From 1,000 Meters

- Three interfaces rolled into one
 - Object Access (like Amazon S3)
 - Block Access
 - (Distributed File System)
- Sitting on top of a Storage Cluster
 - Self Healing
 - Self Managed
 - No Bottlenecks

From 1,000 Meters

Unified Data Handling for 3 Purposes

Object Storage (Like Amazon S3)

- RESTful Interface
- S3 and SWIFT APIs

Block Device

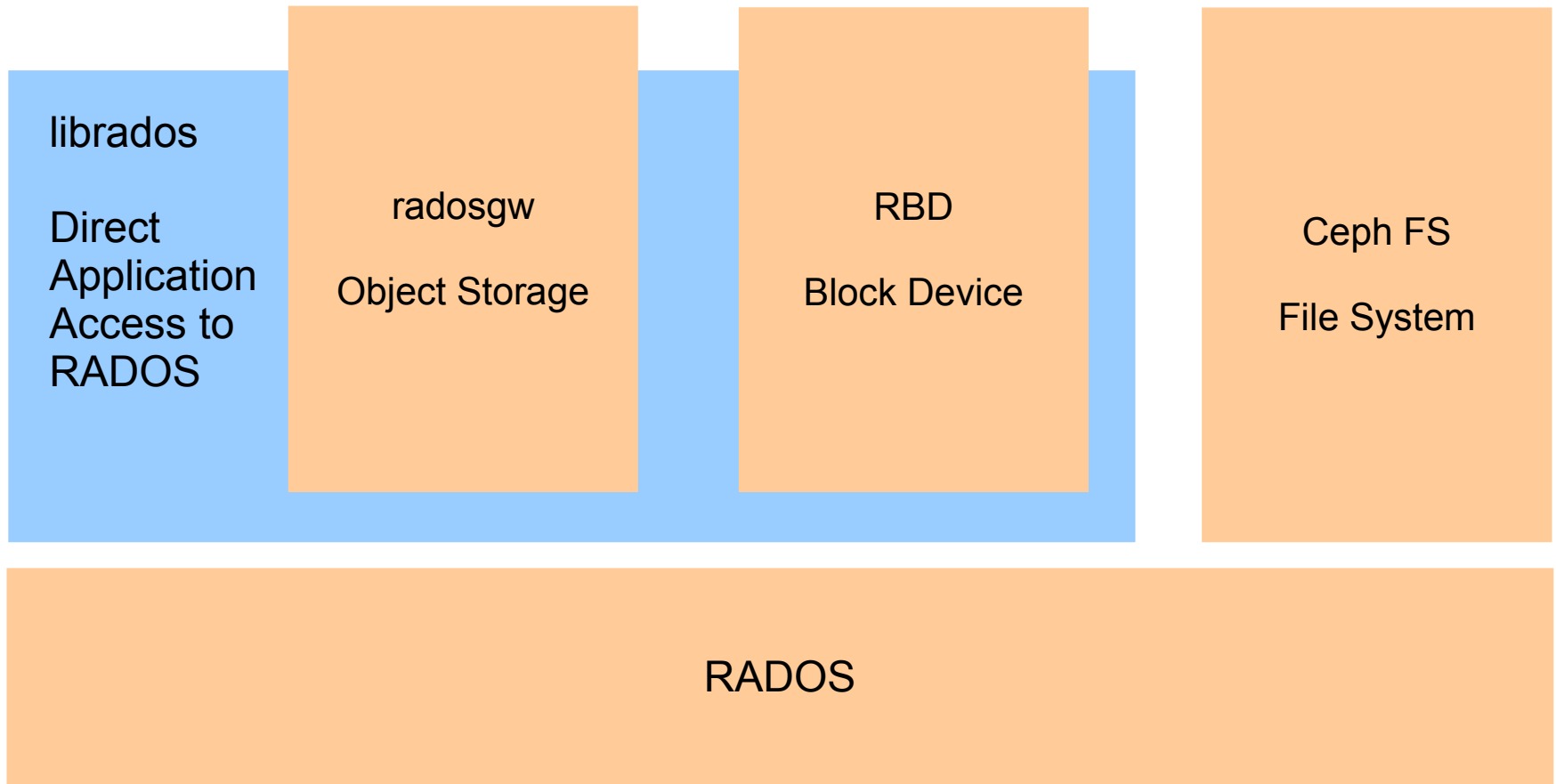
- Block devices
- Up to 16 EiB
- Thin Provisioning
- Snapshots

File System

- POSIX Compliant
- Separate Data and Metadata
- For use e.g. with Hadoop

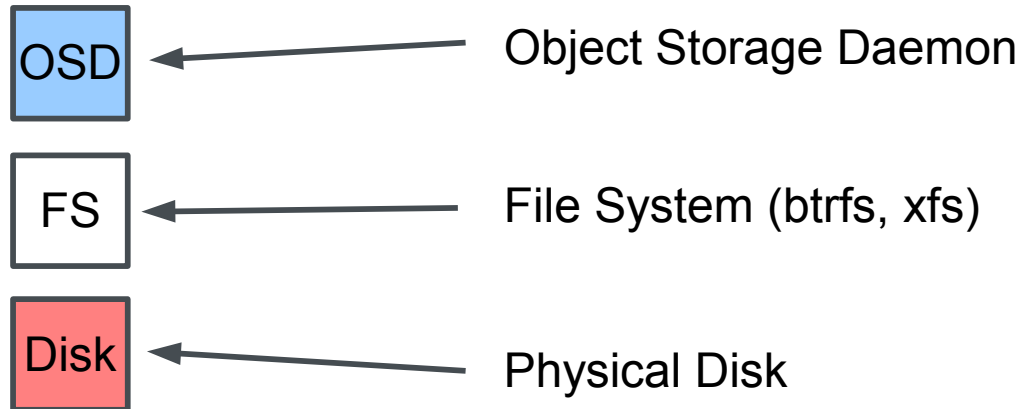
Autonomous, Redundant Storage Cluster

Component Names



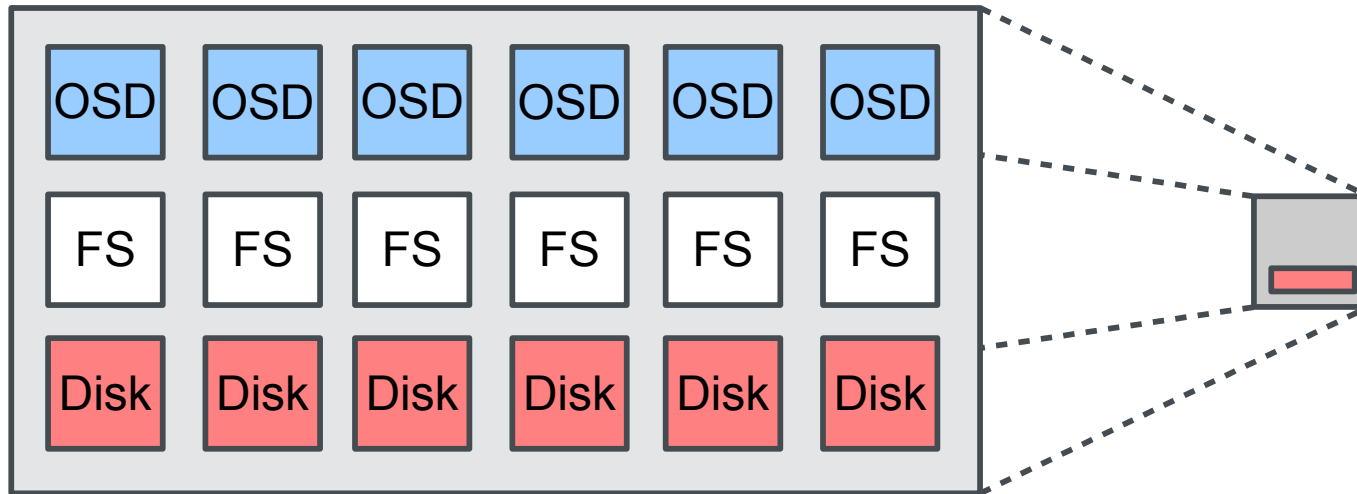
How Does Ceph Work?

For a Moment, Zooming to Atom Level



- OSDs serve storage objects to clients
- Peer to peer to perform replication and recovery

Put Several of These in One Node

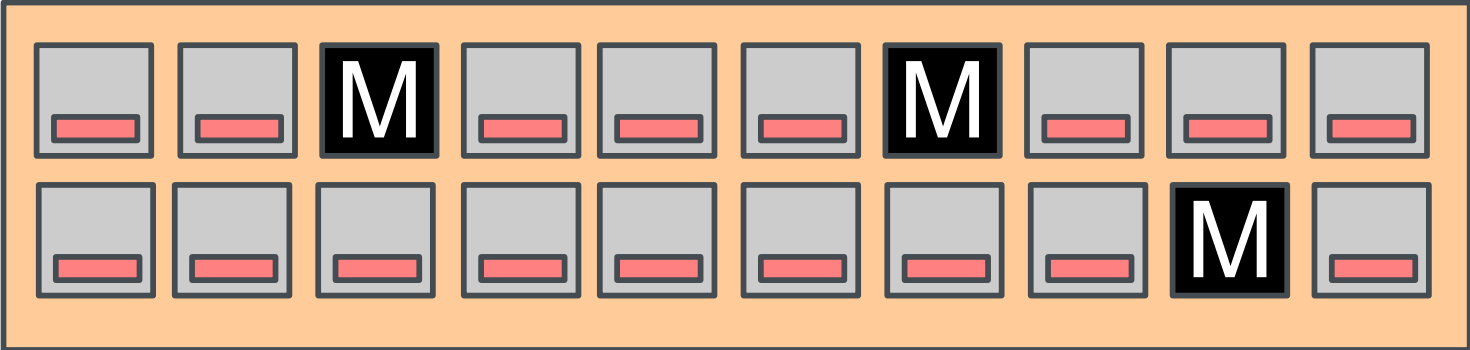


Mix In a Few Monitor Nodes



- Monitors are the brain cells of the cluster
 - Cluster Membership
 - Consensus for Distributed Decision Making
- Do not serve stored objects to clients

Voilà, a Small RADOS Cluster



Several Ingredients

- Basic Idea
 - Coarse grained partitioning of storage supports policy based mapping (don't put all copies of my data in one rack)
 - Topology map and Rules allow clients to “compute” the exact location of any storage object
- Three conceptual components
 - Pools
 - Placement groups
 - CRUSH: deterministic decentralized placement algorithm

Pools

- A pool is a logical container for storage objects
- A pool has a set of parameters
 - a name
 - a numerical ID (internal to RADOS)
 - number of replicas OR erasure encoding settings
 - number of placement groups
 - placement rule set
 - owner
- Pools support certain operations
 - create/remove/read/write entire objects
 - snapshot of the entire pool

Placement Groups

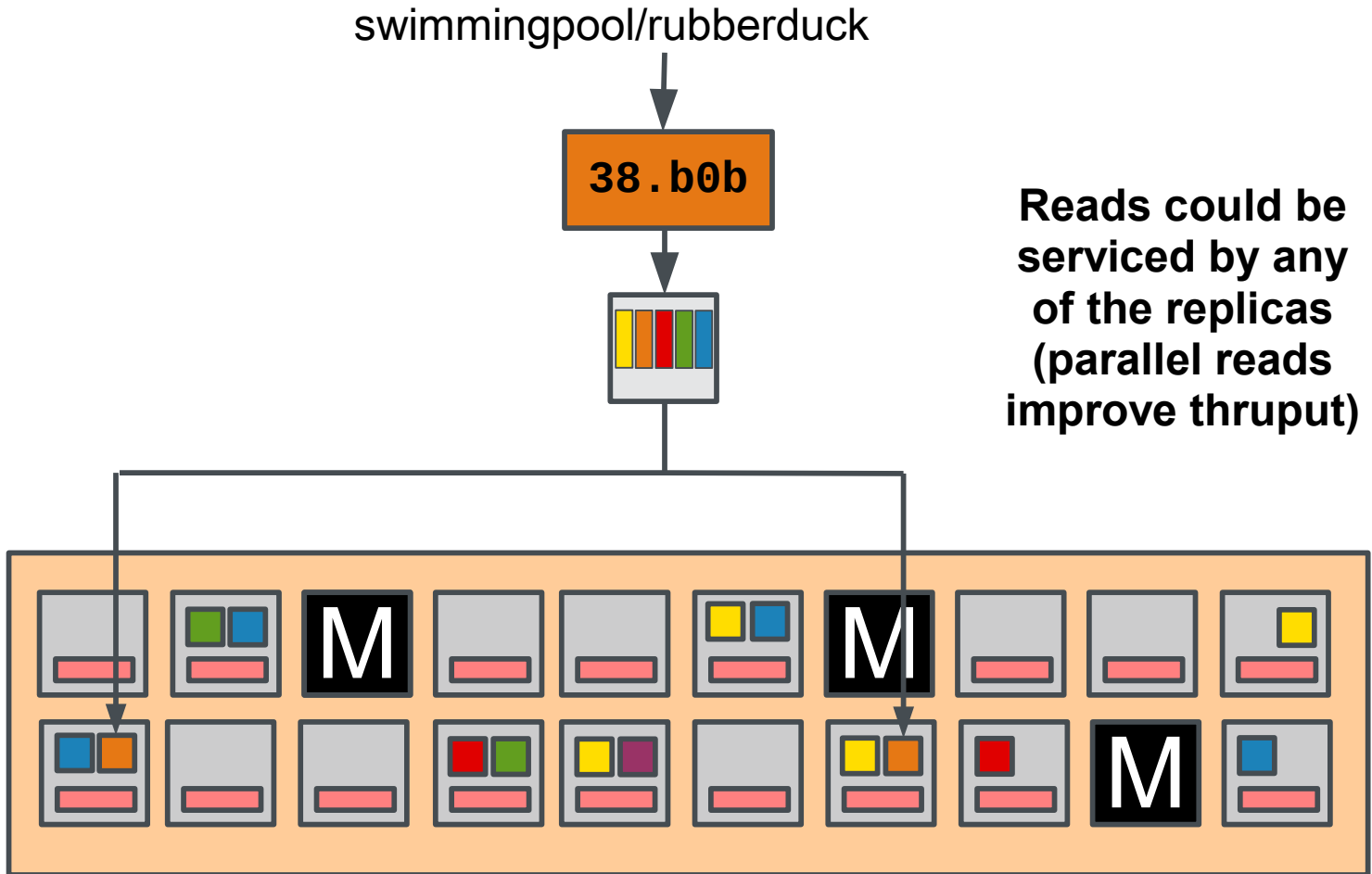
- Placement groups help balance data across OSDs
- Consider a pool named “swimmingpool”
 - with a pool ID of 38 and 8192 placement groups (PGs)
- Consider object “rubberduck” in “swimmingpool”
 - $\text{hash}(\text{“rubberduck”}) \% 8192 = 0xb0b$
 - The resulting PG is 38.b0b
- One PG typically spans several OSDs
 - for balancing
 - for replication
- One OSD typically serves many PGs

CRUSH

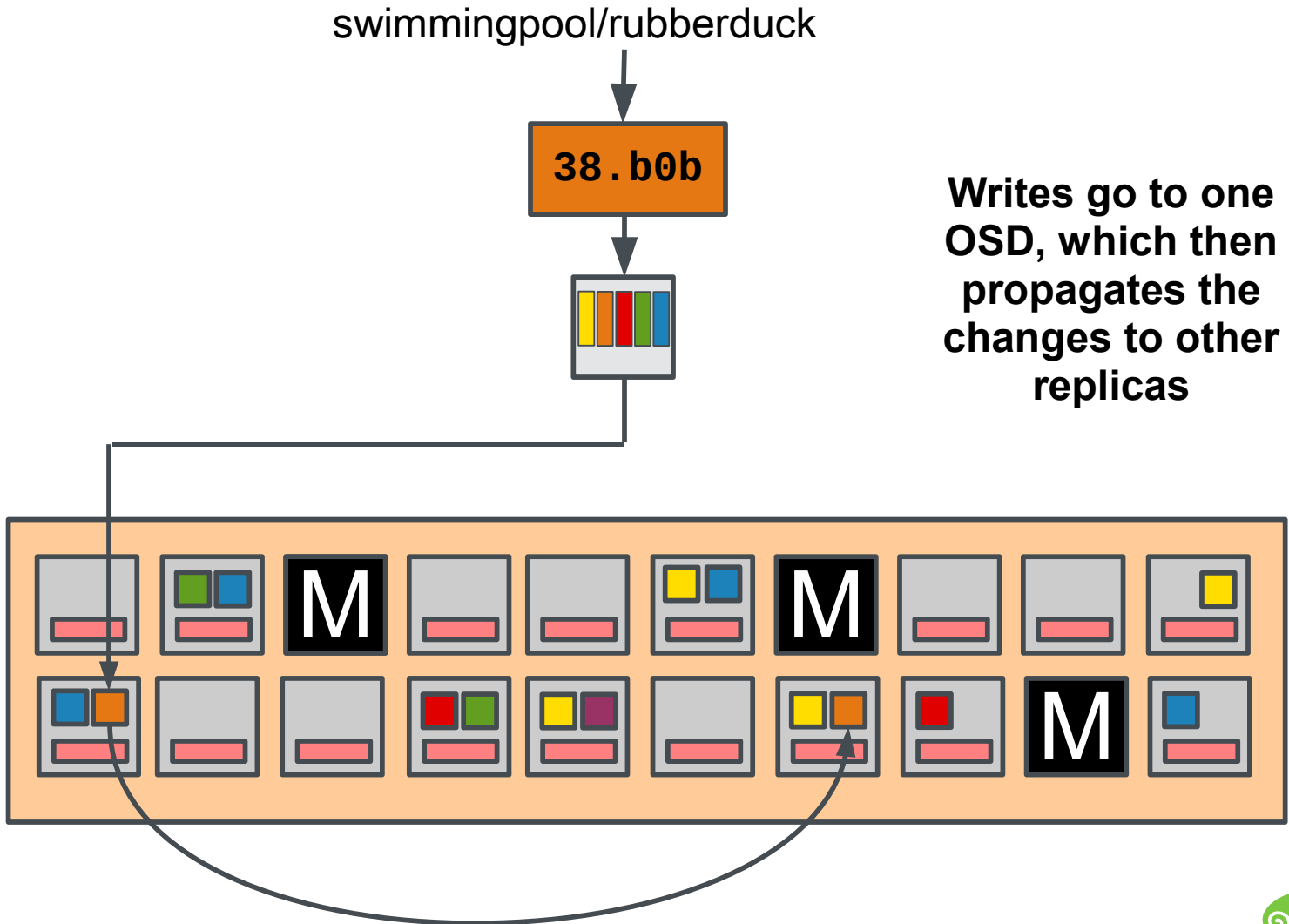


- CRUSH uses a map of all OSDs in your cluster
 - includes physical topology, like row, rack, host
 - includes rules describing which OSDs to consider for what type of pool/PG
- This map is maintained by the monitor nodes
 - Monitor nodes use standard cluster algorithms for consensus building, etc

CRUSH in Action: Reading



CRUSH in Action: Writing



Software Defined Storage

Legacy Storage Arrays

- Limits:

- Tightly controlled environment
- Limited scalability
- Few options
 - Only certain approved drives
 - Constrained number of disk slots
 - Few memory variations
 - Only very few networking choices
 - Typically fixed controller and CPU

- Benefits:

- Reasonably easy to understand
- Long-term experience and “gut instincts”
- Somewhat deterministic behavior and pricing

Software Defined Storage (SDS)

- Limits:

- ?

- Benefits:

- Infinite scalability
 - Infinite adaptability
 - Infinite choices
 - Infinite flexibility

- ... right.

Properties of a SDS System

- Throughput
- Latency
- IOPS

- Availability
- Reliability

- Capacity
- Density

- **Cost**

Architecting a SDS system

- These goals often conflict:
 - Availability versus Density
 - IOPS versus Density
 - Everything versus Cost
- Many hardware options
- Software topology offers many configuration choices
- There is no one size fits all

Setup Choices

Networking (Public and Private)

- Ethernet (1, 10, 40 GbE)
 - Reasonably inexpensive (except for 40 GbE)
 - Can easily be bonded for availability
 - Use jumbo frames
- Infiniband
 - High bandwidth
 - Low latency
 - Typically more expensive
 - No support for RDMA yet in Ceph, need to use IPoIB

Network

- Choose the fastest network you can afford
- Switches should be low latency with fully meshed backplane
- Separate public and cluster network
- Cluster network should typically be twice the public bandwidth
 - Incoming writes are replicated over the cluster network
 - Re-balancing and re-mirroring take utilize the cluster network

Different Access Modes

- radosgw:
 - An additional gateway in front of your RADOS cluster
 - Little impact on throughput, but it does affect latency
- User-space RADOS access:
 - More feature rich than in-kernel rbd.ko module
 - Typically provides higher performance

Storage Node

- CPU
 - Number and speed of cores
- Memory
- Storage controller
 - Bandwidth, performance, cache size
- SSDs for OSD journal
 - SSD to HDD ratio
- HDDs
 - Count, capacity, performance

Adding More Nodes

- Capacity increases
- Total throughput increases
- IOPS increase
- Redundancy increases
- Latency unchanged
- Eventually: network topology limitations
- Temporary impact during re-balancing

Adding More Disks to a Node

- Capacity increases
- Redundancy increases
- Throughput might increase
- IOPS might increase
- Internal node bandwidth is consumed
- Higher CPU and memory load
- Cache contention
- Latency unchanged

OSD File System

- btrfs

- Typically better write throughput performance
- Higher CPU utilization
- Feature rich
 - Compression, checksums, copy on write
- The choice for the future!

- XFS

- Good all around choice
- Very mature for data partitions
- Typically lower CPU utilization
- The choice for today!

Impact of Caches

- Cache on the client side
 - Typically, biggest impact on performance
 - Does not help with write performance
- Server OS cache
 - Low impact: reads have already been cached on the client
 - Still, helps with readahead
- Caching controller, battery backed:
 - Significant impact for writes

Impact of SSD Journals

- SSD journals accelerate bursts and random write IO
- For sustained writes that overflow the journal, performance degrades to HDD levels
- SSDs help very little with read performance
- SSDs are very costly
 - ... and consume storage slots -> lower density
- A large battery-backed cache on the storage controller is highly recommended if not using SSD journals

Hard Disk Parameters

- Capacity matters
 - Often, highest density is not most cost effective
- On-disk cache matters less
- Reliability advantage of Enterprise drives typically marginal compared to cost
 - Buy more drives instead
- RPM:
 - Increase IOPS & throughput
 - Increases power consumption
 - 15k drives quite expensive still

Impact of Redundancy Choices

- Replication:

- n number of exact, full-size copies
- Potentially increased read performance due to striping
- Increased cluster network utilization for writes
- Rebuilds can leverage multiple sources
- Significant capacity impact

- Erasure coding:

- Data split into k parts plus m redundancy codes
- Better space efficiency
- Higher CPU overhead
- Significant CPU and cluster network impact, especially during rebuild
- Cannot directly be used to with block devices (see next slide)

Cache Tiering

- Multi-tier storage architecture:
 - Pool acts as a transparent write-back overlay for another
 - e.g., SSD 3-way replication over HDDs with erasure coding
 - Can flush either on relative or absolute dirty levels, or age
 - Additional configuration complexity and requires workload-specific tuning
 - Also available: read-only mode (no write acceleration)
 - Some downsides (no snapshots)
- A good way to combine the advantages of replication and erasure coding

Conclusion

Questions and Answers?

Thank you.





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